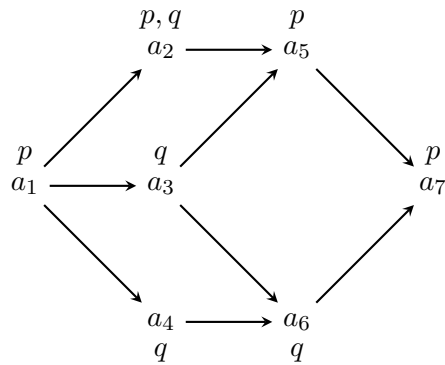




1. Consider the model  $\mathcal{M} = (W, R, V)$  given by the following picture:



- (a) Write out the definitions of  $W$ ,  $R$  and  $V$ .
- (b) Show the following:
  - (i)  $\mathcal{M}, a_1 \models \Box\Box(p \vee q)$ ,
  - (ii)  $\mathcal{M}, a_2 \models \Diamond q \rightarrow \Diamond\Diamond q$ ,
  - (iii)  $\mathcal{M}, a_3 \models \Diamond p \rightarrow \Box(q \rightarrow \Box(p \rightarrow \Box p))$ .
- (c) Show the following:
  - (i)  $\mathcal{M} \not\models p \rightarrow \Diamond p$ ,
  - (ii)  $\mathcal{M} \models \Box\Box\Box\neg q$ ,
  - (iii)  $\mathcal{M} \not\models q \rightarrow (\Diamond q \rightarrow \Box(q \rightarrow \Diamond q))$ .
- (d) Change the valuation  $V$  on the frame such that in the new model  $\mathcal{M}' = (W, R, V')$  it holds that:  $\mathcal{M}' \models \Box p \rightarrow p$ .

2. The *binary tree* is the frame  $\mathcal{B} = (W, R)$  where the domain  $W$  is the set of all finite strings over the alphabet  $\{0,1\}$ :

$$W = \{0,1\}^* = \{\varepsilon, 0, 1, 00, 01, 10, 11, 000, \dots\}.$$

Here  $\varepsilon$  denotes the empty string, and the transition relation  $R \subseteq W \times W$  is given by:

$$Rst \quad \text{if and only if} \quad t = s0 \text{ or } t = s1$$

- (a) Make a drawing of the first four levels of  $\mathcal{B}$ .
- (b) Consider a valuation  $V$  on  $\mathcal{B}$  that makes  $p$  true on all strings of even length. Show that  $\mathcal{B}, V \models \Box \Diamond p \rightarrow \Diamond \Box p$ .
- (c) Let  $V'$  be a valuation on  $\mathcal{B}$  which makes the variable  $p$  true on all strings whose first letter is 0, and  $q$  on strings with first letter 1, so:

$$V'(p) = \{0w \mid w \in \{0,1\}^*\}$$

$$V'(q) = \{1w \mid w \in \{0,1\}^*\}$$

Use this valuation to show that the formula  $\lambda$  defined by:

$$\lambda = \Diamond p \wedge \Diamond q \rightarrow \Diamond(p \wedge \Diamond q) \vee \Diamond(p \wedge q) \vee \Diamond(\Diamond p \wedge q)$$

is not valid in the binary tree.

- (d) Show that the formula  $\Diamond \Diamond p \rightarrow \Diamond p$  is not valid in  $\mathcal{B}$ .

3. Prove or disprove universal validity of the following formulas:

- (a)  $\Box(p \rightarrow q) \rightarrow (\Diamond p \rightarrow \Diamond q)$
- (b)  $\Box(p \wedge q) \rightarrow (\Diamond p \wedge \Diamond q)$
- (c)  $\Box p \rightarrow \Diamond p$
- (d)  $\Diamond p \rightarrow \Box p$
- (e)  $\Box(\Box p \rightarrow p) \rightarrow \Box p$